

Program Semantics in Model-Based WCET Analysis

— State of the Art Perspective —

Mihail Asavoaie

Claire Maiza

Pascal Raymond

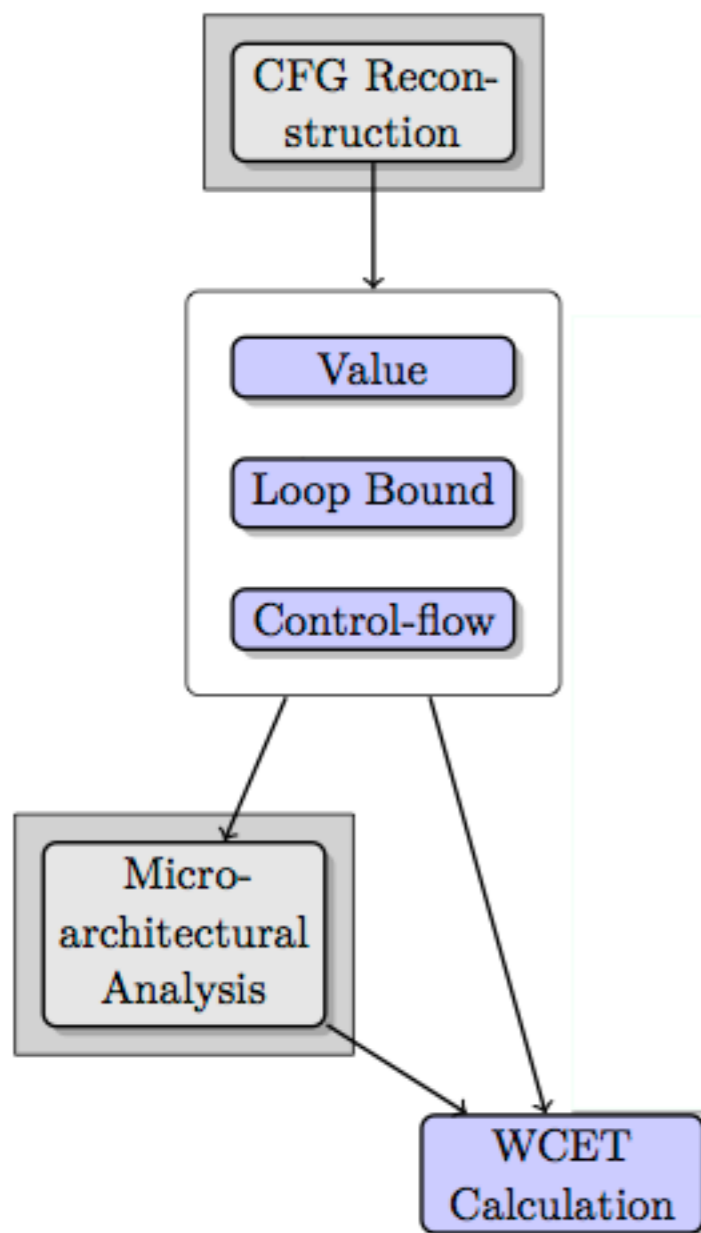


W-SEPT

WCET: **S**emantics, **P**recision and **T**raceability

WCET 2013 - PARIS

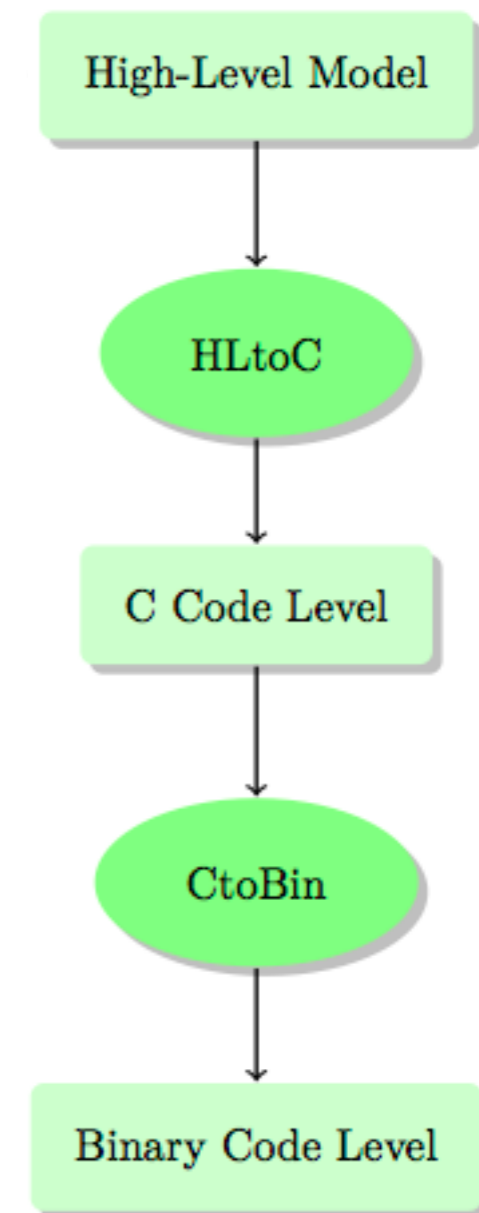
Key Point! - Advocate for WCET analysis based on the program semantics



(1) Model-Based Design

+

(2) WCET Analysis



- Synchronous Programming

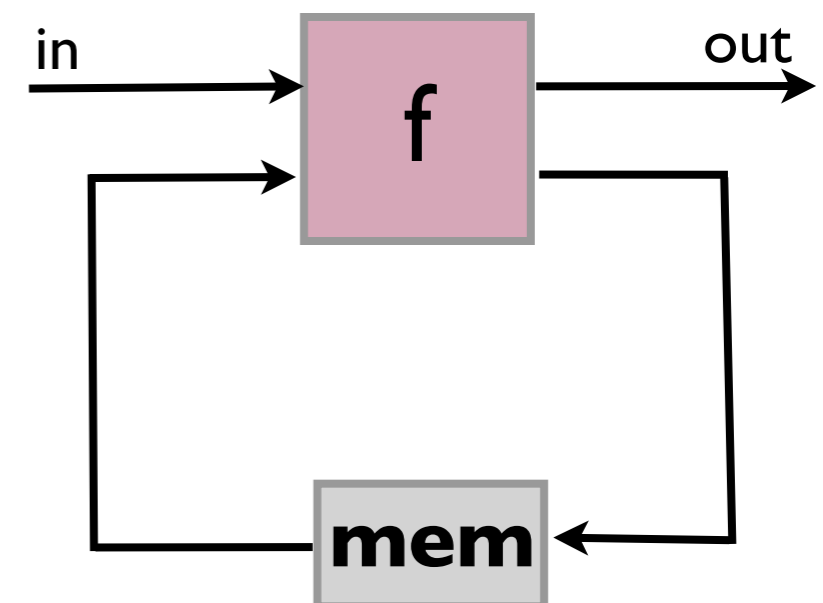
“This spirit is in full accordance with modern thinking about system design, sometimes called **model-based design**, which advocates the use of high-level models with **“ideal”, implementation-independent semantics**, and **the separation of concerns between design and implementation.**”

“Synchronous Programming” - P. Caspi et al.
in Handbook of Real-Time and Embedded Systems - 2008

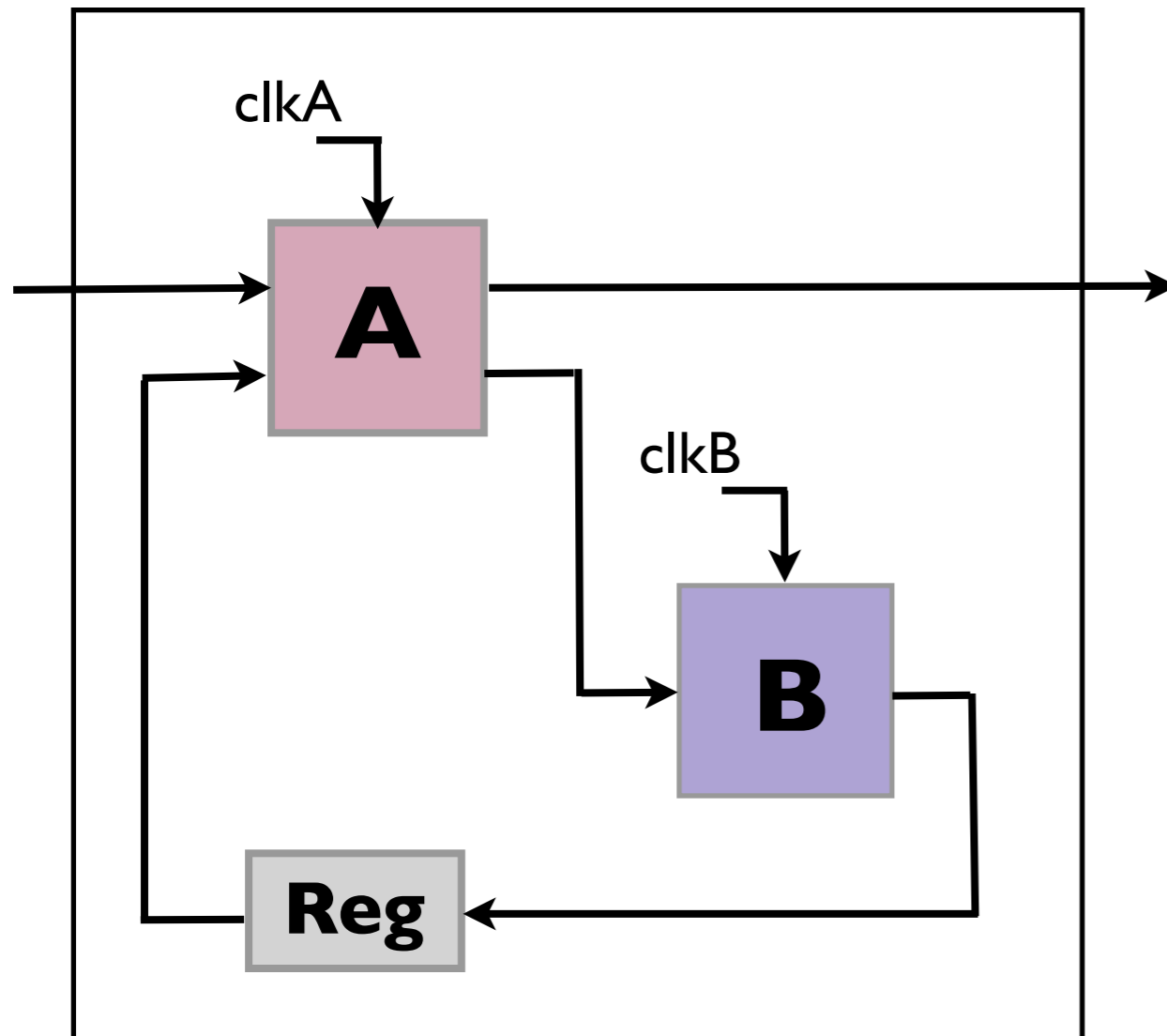
- Generic programming paradigm - **designed to generate code**
- deterministic concurrent programming style
- in **data-flow** (e.g. Lustre), **control-flow** (e.g. Esterel) flavors, or **both** (e.g. Scade 6 - industrial version of Lustre + finite state machines)

- Lustre

- **semantics** - functions over flows
- **style** - network of operating nodes (similar to sequential circuits or block diagrams)
- Scade - in critical systems: transportation, energy



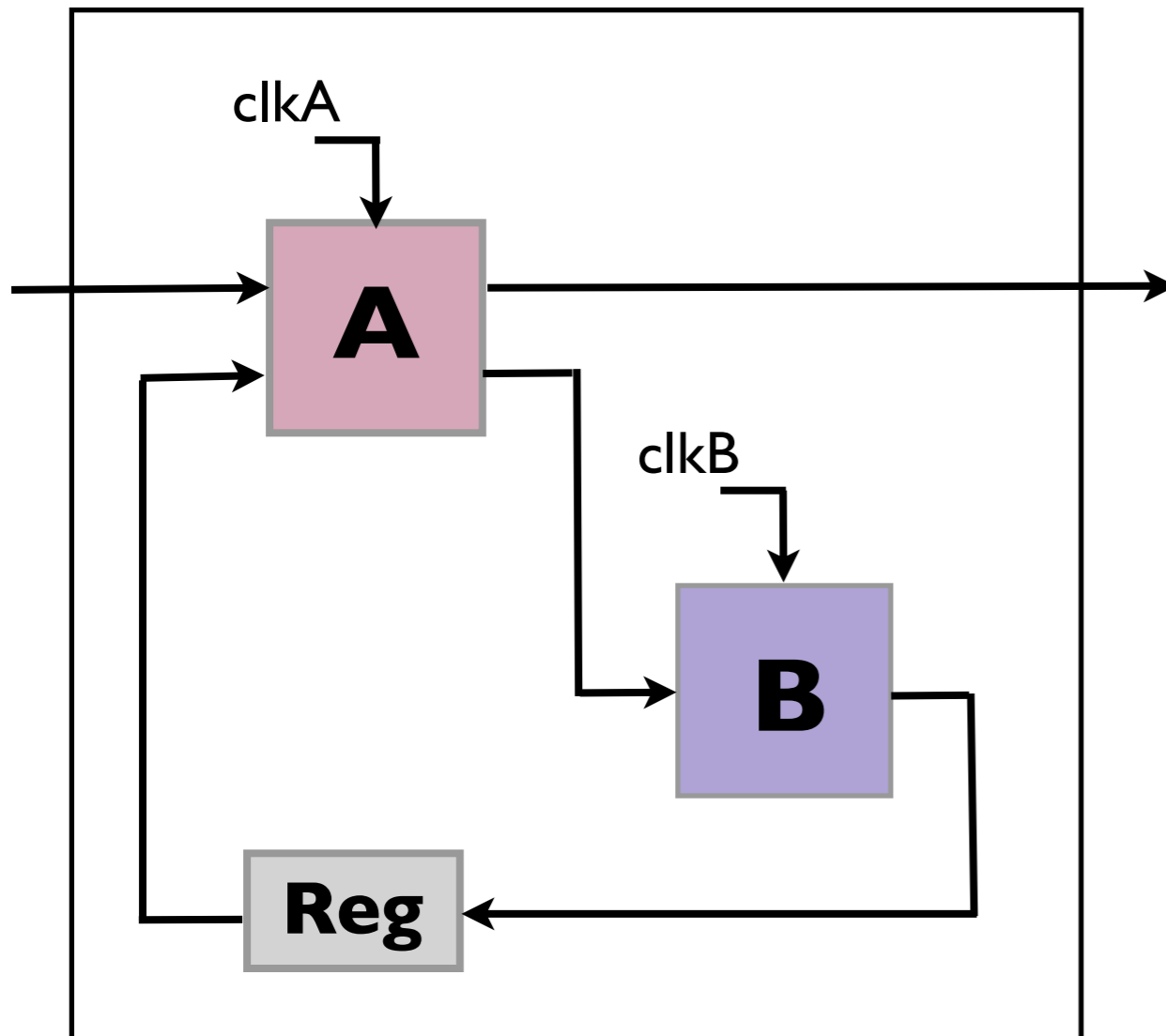
Design



Elements

- overall → subsystems (**A**, **B**)
- subsystem → inter-dependent data flows
- dependencies → direct and logical (delays, i.e. **Reg**)
- clocks (*clkA* and *clkB*)

Design



Elements

→ “... *the separation of concerns between design and implementation.*”

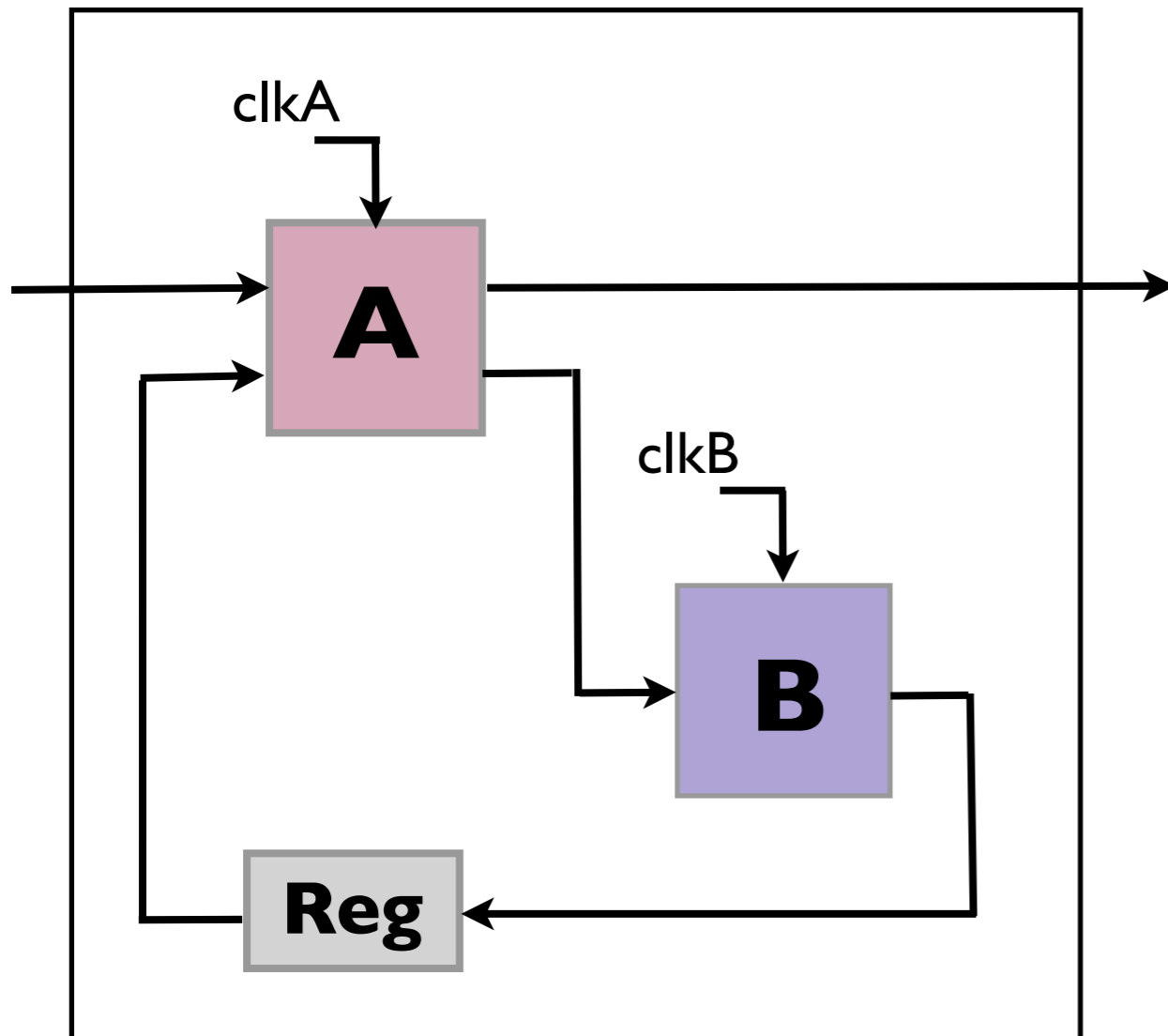
Design:

- instantaneous execution

Implementation:

- bounded execution (with a ***guaranteed upper bound on time***)

Design



C Code

- scheduling → **A** before **B**

- code →

```
global Reg;
```

```
if clkA then A_Step()
```

```
if clkB then {
```

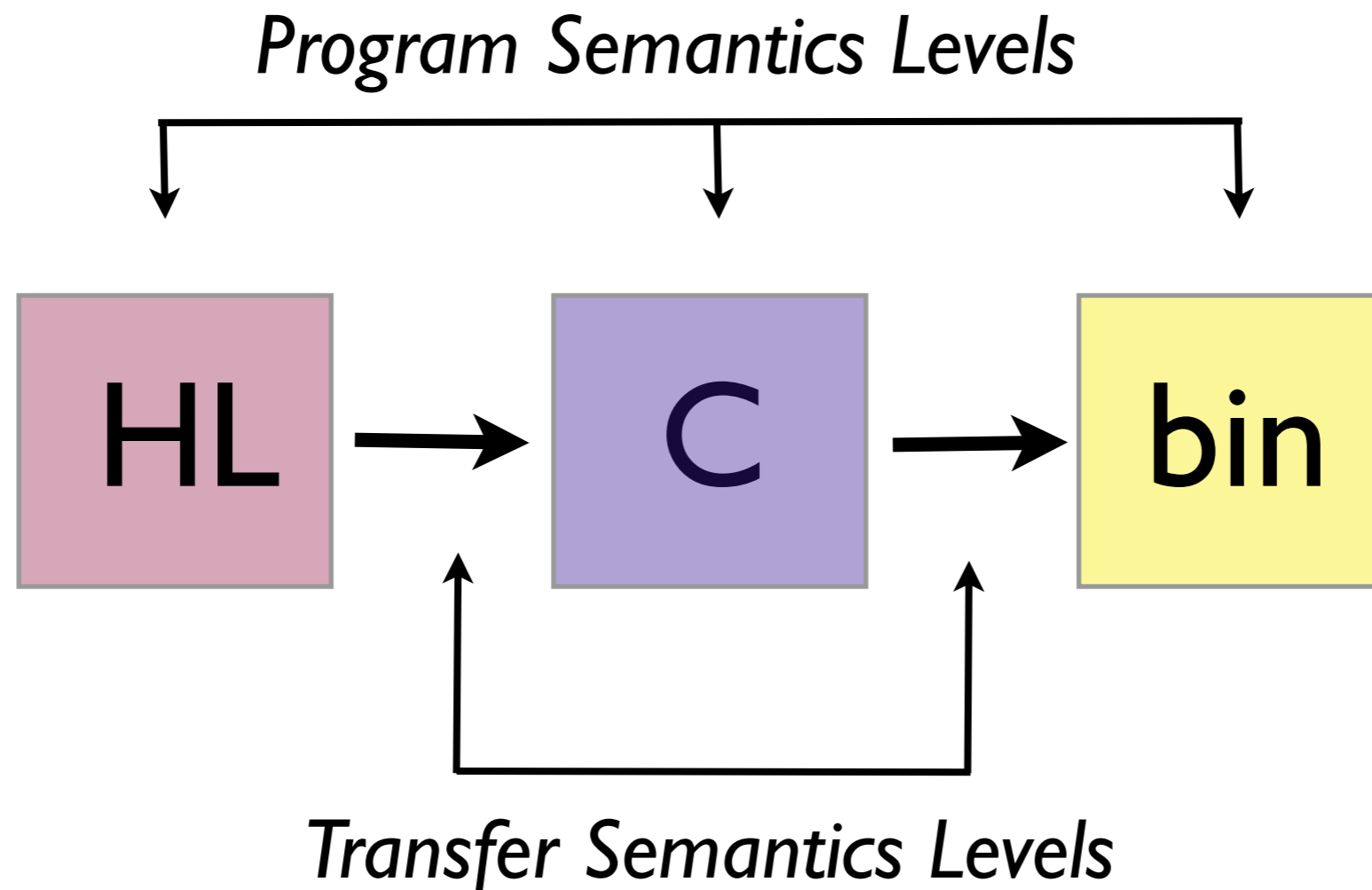
```
    B_Step();
```

```
    Reg = B.out;
```

```
}
```

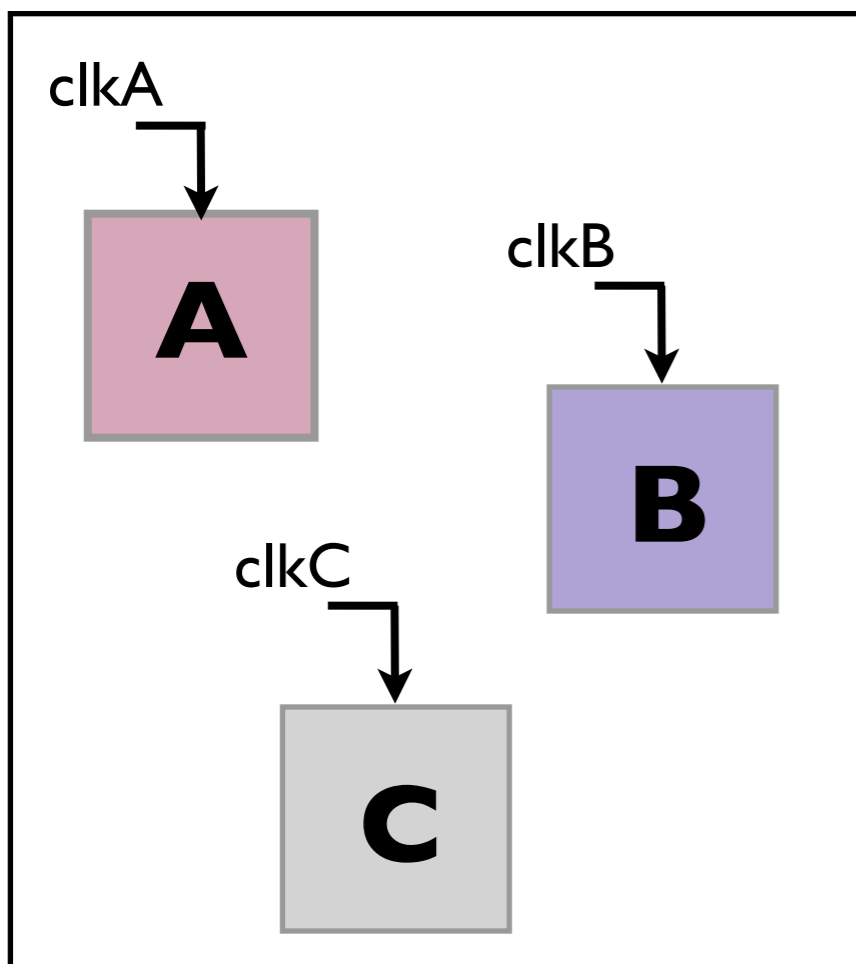
- Deterministic, traceable C code (heap-free, no recursion, bounded loops)

- Typical Synchronous Program + Compilation



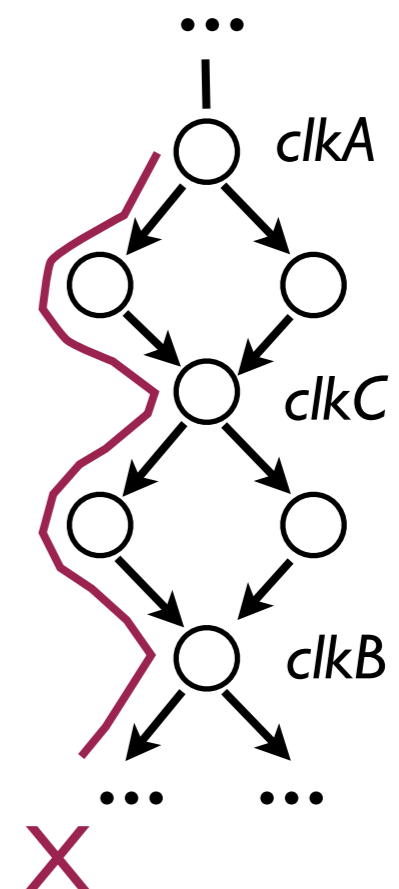
- From the WCET analysis perspective : both global and local interest

- Generated code from model-based designs:
 - subsystems
 - inter-dependent
 - **clocks (Lustre)**

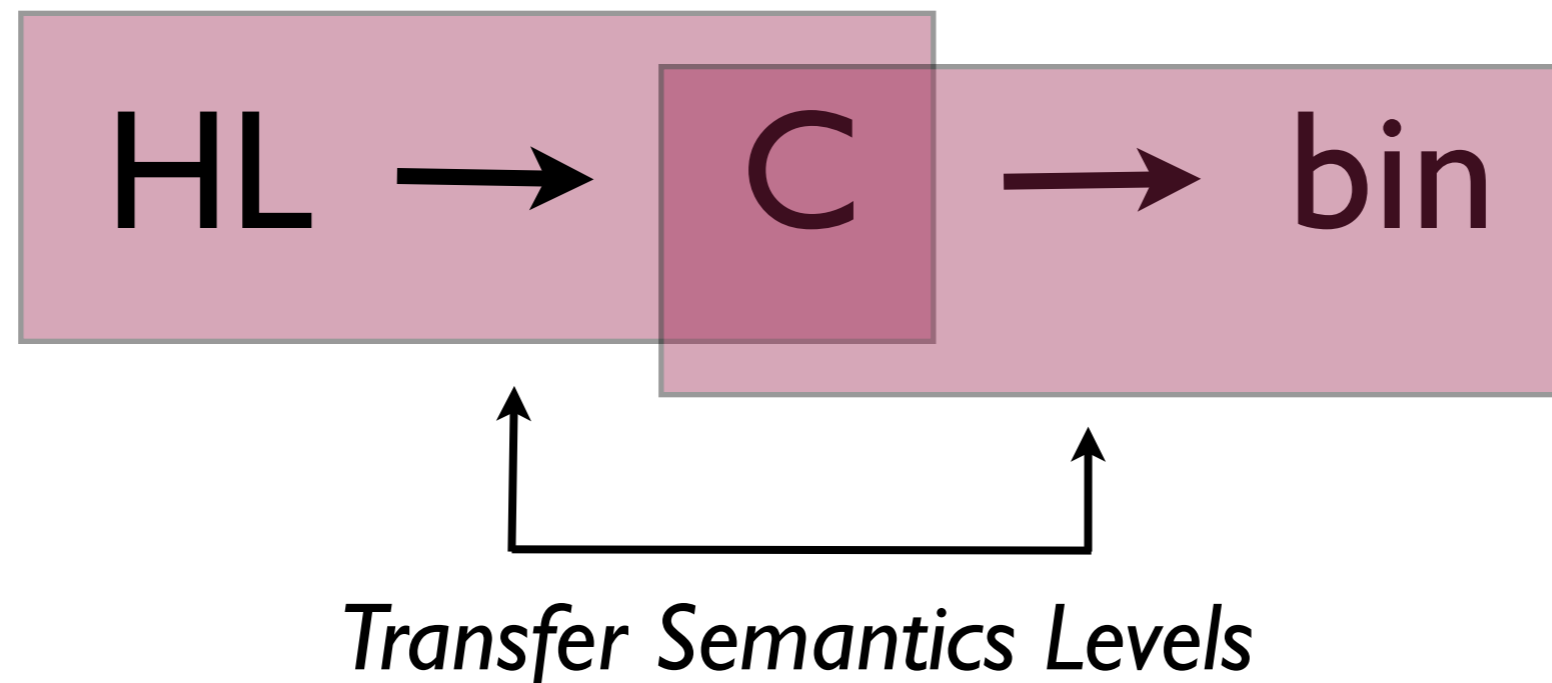


- assume a relation on the nodes: at most one of **A**, **B**, **C** is not executed

scheduling analysis

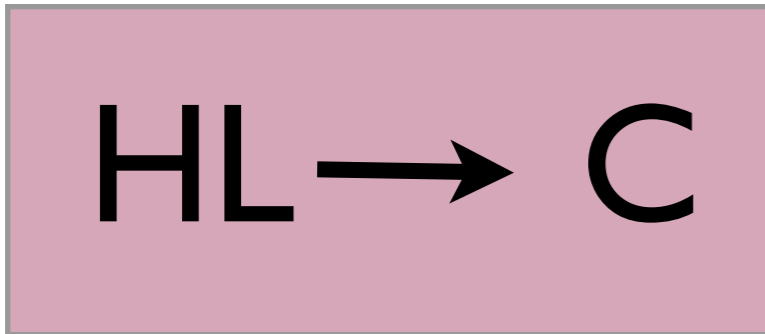


- Two levels of interest:



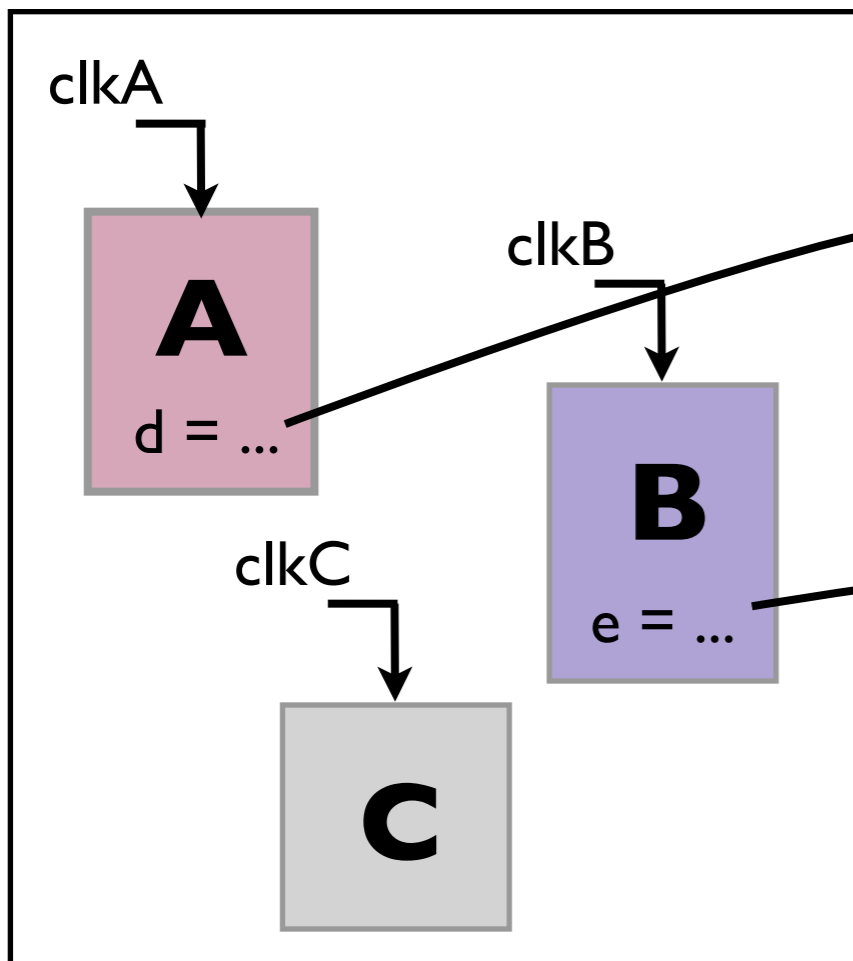
- **HL to C** : it features a higher degree of control over the model semantics transformation
- **C to bin** : compilation with optimizations

- identify relations between different representations



.C FILE

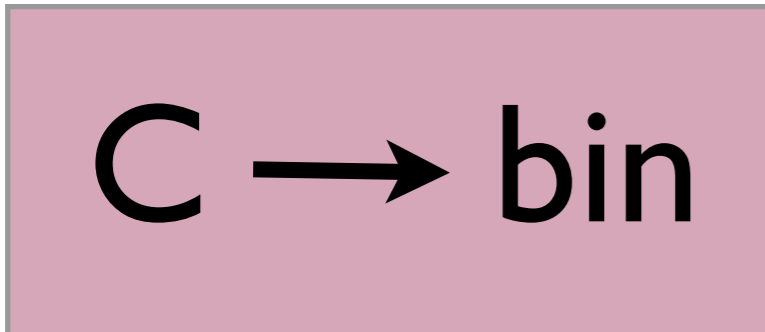
```
if clkA then A_Step()
if clkB then B_Step()
```



```
A_Step() {
  d = ...
}
```

```
B_step() {
  e = ...
}
```

- transfer of semantics through optimizations



```
if clkA then A_Step()  
if clkB then B_Step()
```

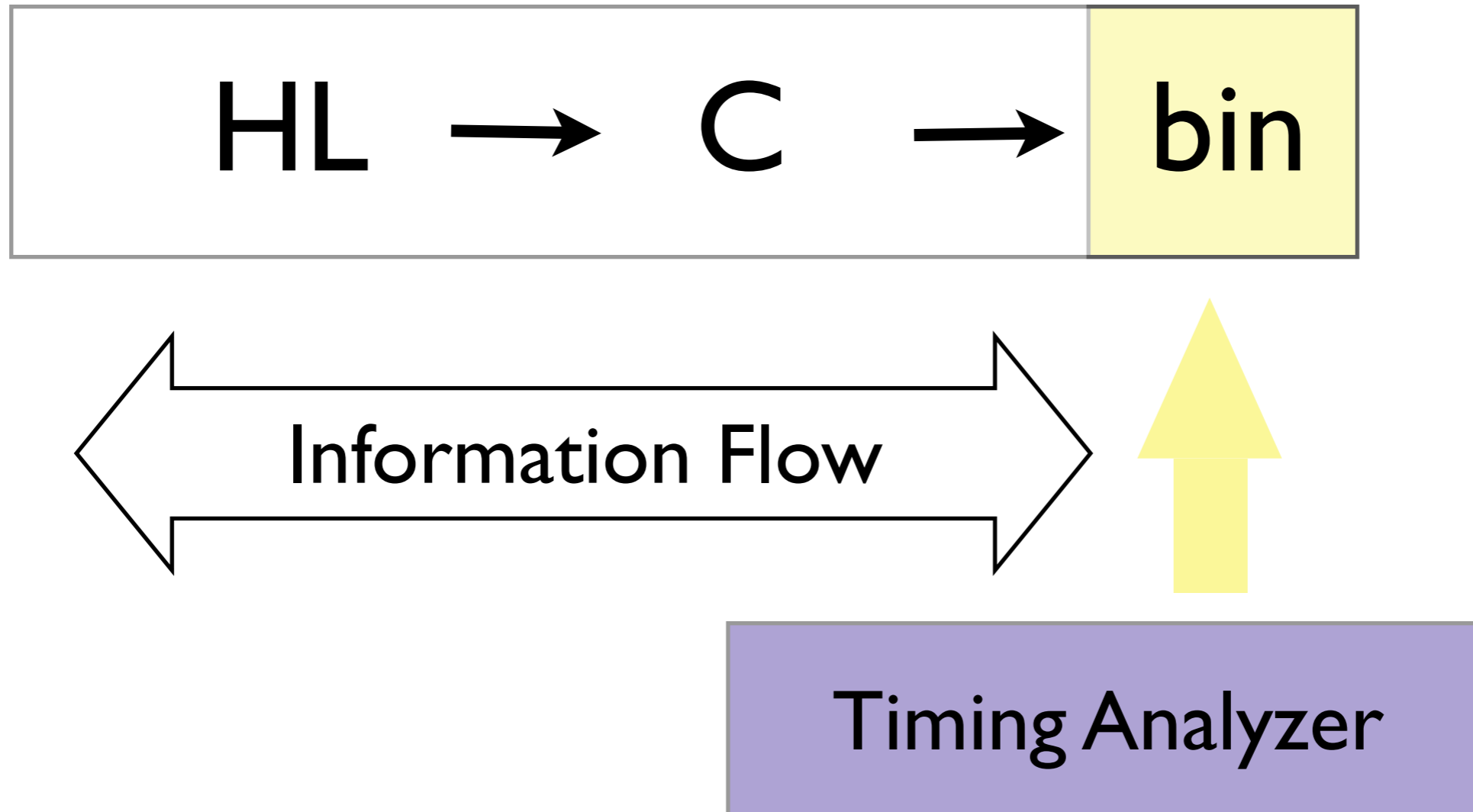
```
A_Step() {  
    d = ...  
}  
B_step() {  
    e = ...  
}
```

- **obfuscation of the control structure**

- no clear solution, in general optimizations are handled in a case by case manner

- specialized annotation language
- transfer functions

- Integration methods

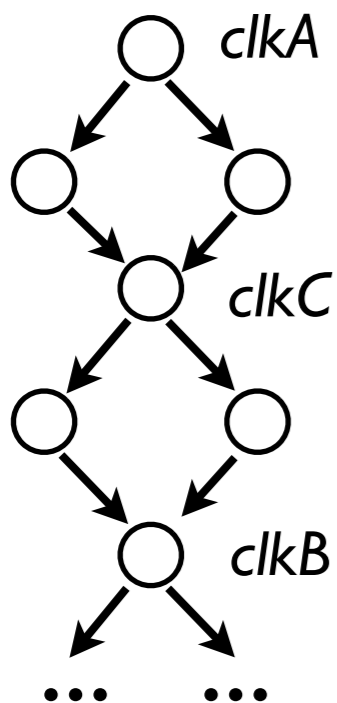


- seamless integration vs modified MBD infrastructure

- the path analysis determines the worst-case execution time based on previously computed (or annotated) artifacts

- represented outside MBD

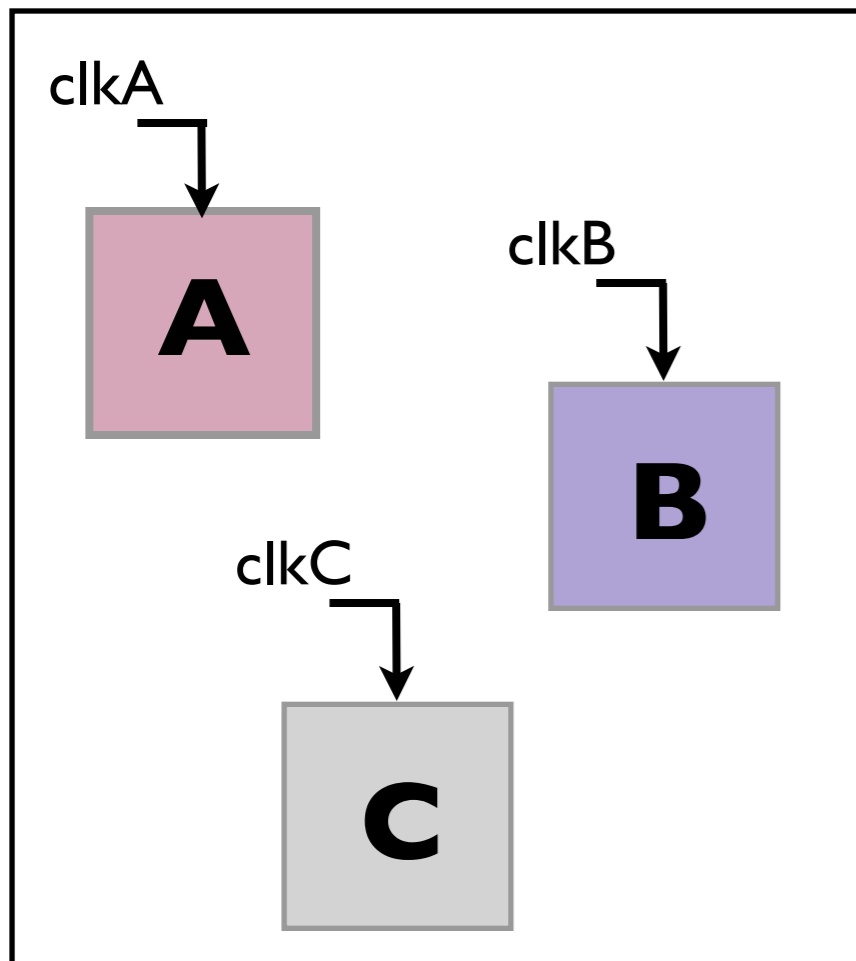
- **IPET** - alternatives: on the program syntax, on model checking



Is IPET the best approach to represent the semantics?

- separation of concerns at the level of WCET analysis workflow

- IPET - path analysis through ILP constraint solving
 - at most one of **A**, **B**, **C** is not executed

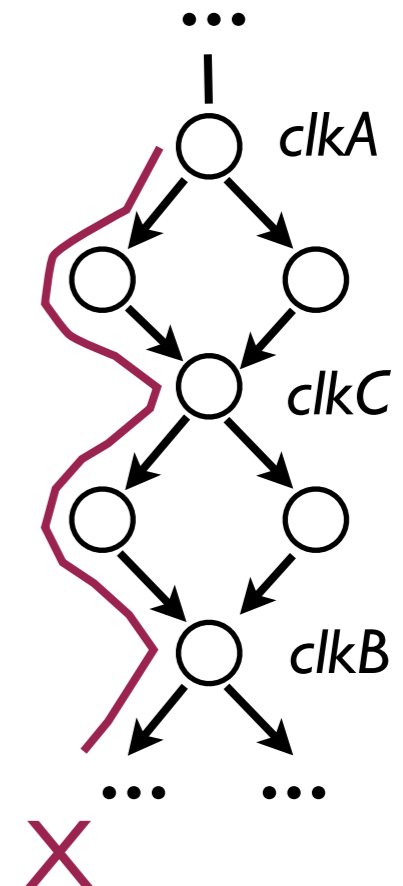


Solve:

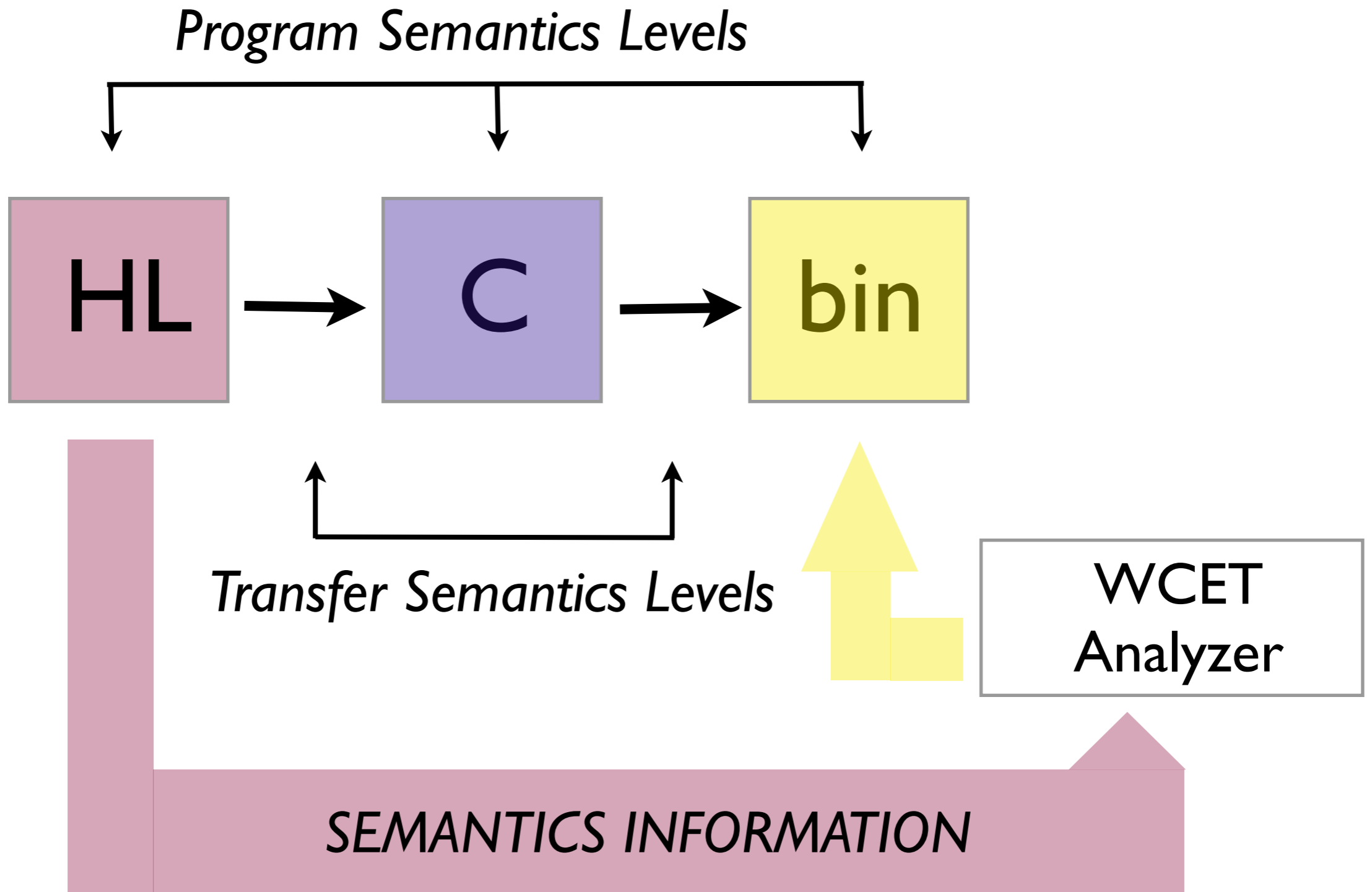
$$X_{\mathbf{A}} + X_{\mathbf{B}} + X_{\mathbf{C}} \leq 2$$

and

ILP representation of the control structure



- IPET enhancements are possible



- the **paper** presents a state-of-the-art overview of approaches for both local and global w.r.t. MBD workflow, from the WCET analysis perspective
 - particular MBD: synchronous programs as models and C code generation capabilities
- the **presentation** (complements the paper)
 - follows several aspects of the synchronous programming
 - advocates for a separation of concerns at the level of WCET analysis

